

## **Learning and Computing**

Data Mining (データマイニング)

Mahito Sugiyama (杉山麿人)

## **Objective of Today's Lecture**

- Learn a fundamental mechanism of machine learning
  - Machine learning is a core process in many applications in data mining
- Computational aspects of machine learning are mainly discussed
- Key issues:
  - Computing (single) vs Learning (double)
    - Finite/infinite
  - Learning targets (mathematical objects) vs
     Representations (programs)

[Schoelkopf, 2013]

[Schoelkopf, 2013]

• 1, 2, 4, 7, ...  $\rightarrow$  What are succeeding numbers? 1, 2, 4, 7, 11, 16, ...  $(a_n = a_{n-1} + n - 1)$ 

#### [Schoelkopf, 2013]

1, 2, 4, 7, 11, 16, ... 
$$(a_n = a_{n-1} + n - 1)$$
  
1, 2, 4, 7, 12, 20, ...  $(a_n = a_{n-1} + a_{n-2} + 1)$ 

#### [Schoelkopf, 2013]

1, 2, 4, 7, 11, 16, ... 
$$(a_n = a_{n-1} + n - 1)$$
  
1, 2, 4, 7, 12, 20, ...  $(a_n = a_{n-1} + a_{n-2} + 1)$   
1, 2, 4, 7, 13, 24, ...  $(a_n = a_{n-1} + a_{n-2} + a_{n-3})$ 

#### [Schoelkopf, 2013]

```
1, 2, 4, 7, 11, 16, ... (a_n = a_{n-1} + n - 1)

1, 2, 4, 7, 12, 20, ... (a_n = a_{n-1} + a_{n-2} + 1)

1, 2, 4, 7, 13, 24, ... (a_n = a_{n-1} + a_{n-2} + a_{n-3})

1, 2, 4, 7, 14, 28 (divisors of 28)
```

#### [Schoelkopf, 2013]

```
1, 2, 4, 7, 11, 16, ... (a_n = a_{n-1} + n - 1)

1, 2, 4, 7, 12, 20, ... (a_n = a_{n-1} + a_{n-2} + 1)

1, 2, 4, 7, 13, 24, ... (a_n = a_{n-1} + a_{n-2} + a_{n-3})

1, 2, 4, 7, 14, 28 (divisors of 28)

1, 2, 4, 7, 1, 1, 5, ... (\pi = 3.1415 \dots \text{ and } e = 2.718 \dots)
```

#### [Schoelkopf, 2013]

• 1, 2, 4, 7, ...  $\rightarrow$  What are succeeding numbers?

```
1, 2, 4, 7, 11, 16, ... (a_n = a_{n-1} + n - 1)

1, 2, 4, 7, 12, 20, ... (a_n = a_{n-1} + a_{n-2} + 1)

1, 2, 4, 7, 13, 24, ... (a_n = a_{n-1} + a_{n-2} + a_{n-3})

1, 2, 4, 7, 14, 28 (divisors of 28)

1, 2, 4, 7, 1, 1, 5, ... (\pi = 3.1415 \dots \text{ and } e = 2.718 \dots)
```

More than 1229 rules! (https://oeis.org)

#### **Analyze Learning as Scientific Problem**

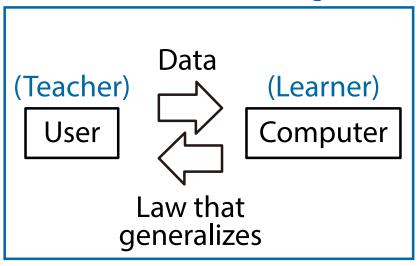
- Which is the correct answer (or generalization)
   for succeeding numbers of 1, 2, 4, 7, ... ?
  - Any answer is possible!

#### **Analyze Learning as Scientific Problem**

- Which is the correct answer (or generalization) for succeeding numbers of 1, 2, 4, 7, ... ?
  - Any answer is possible!
- We should take two points into consideration:
  - (i) We need to formalize the problem of "learning"
    - There are two agents (teacher and learner) in learning, which are different from "computation"
  - (ii) Learning is an infinite process
    - A learner usually never knows that the current hypothesis is perfectly correct

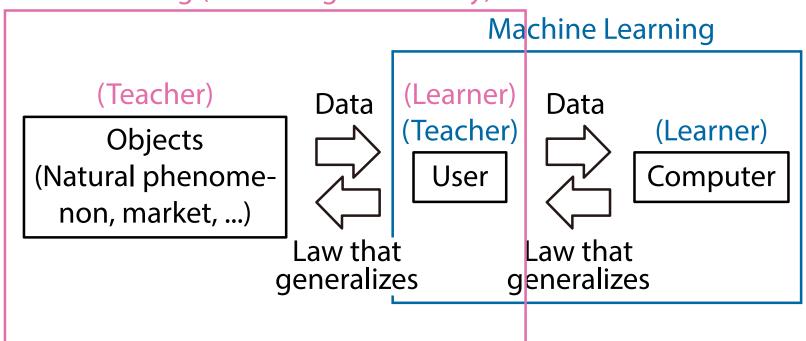
#### Framework of Learning (ML vs DM)

#### Machine Learning



#### Framework of Learning (ML vs DM)

Data Mining (Knowledge Discovery)



#### **Computation** — Core Engine of Learning/Mining

- Machine learning/data mining is usually achieved using a computer
- Computing behavior is mathematically formulated by Alan Turing in 1936
  - A. M. Turing, On Computable Numbers, with the Application to the Entscheidungsproblem, Proceedings of the London Mathematical Society, 42(1), 230–265, 1937
- The model of computation, known as a Turing machine, is developed for simulating computation by human beings

## ON COMPUTABLE NUMBERS, WITH AN APPLICATION TO THE ENTSCHEIDUNGSPROBLEM

By A. M. TURING.

[Received 28 May, 1936.—Read 12 November, 1936.]

The "computable" numbers may be described briefly as the real numbers whose expressions as a decimal are calculable by finite means. Although the subject of this paper is ostensibly the computable numbers.

8/34

Computing is normally done by writing certain symbols on paper. may suppose this paper is divided into squares like a child's arithmetic book. In elementary arithmetic the two-dimensional character of the paper is sometimes used. But such a use is always avoidable, and I think that it will be agreed that the two-dimensional character of paper is no essential of computation. I assume then that the computation is carried out on one-dimensional paper, i.e. on a tape divided into squares. I shall also suppose that the number of symbols which may be printed is finite. If we were to allow an infinity of symbols, then there would be symbols differing to an arbitrarily small extent. The effect of this restriction of the number of symbols is not very serious. It is always possible to use sequences of symbols in the place of single symbols. Thus an Arabic numeral such as

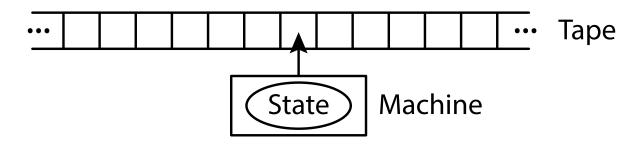
<sup>†</sup> If we regard a symbol as literally printed on a square we may suppose that the square is  $0 \le x \le 1$ ,  $0 \le y \le 1$ . The symbol is defined as a set of points in this square, viz. the set occupied by printer's ink. If these sets are restricted to be measurable, we can define the "distance" between two symbols as the cost of transforming one symbol into the other if the cost of moving unit area of printer's ink unit distance is unity, and there is an infinite supply of ink at x = 2, y = 0. With this topology the symbols form a conditionally compact space.

The behaviour of the computer at any moment is determined by the symbols which he is observing, and his "state of mind" at that moment.

We may suppose that there is a bound B to the number of symbols or squares which the computer can observe at one moment. If he wishes to observe more, he must use successive observations. We will also suppose that the number of states of mind which need be taken into account is finite. The reasons for this are of the same character as those which restrict the number of symbols. If we admitted an infinity of states of mind, some of

them will be "arbitrarily close" and will be confused. Again, the restriction

## **Turing Machine**



- The machine repeats the following:
  - Read a symbol *a* of a cell
  - Do the following from the symbol a and the current state s according to a set of rules in its memory
    - Replace the symbol a at the square
    - Move the head
    - Change the state s

## **Computing vs Learning**

- In *computation*, the process is completed on its own
  - No interaction
    - The Turing machine automatically works according to programmed rules
  - A finite process

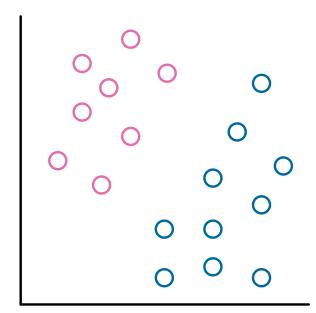
## **Computing vs Learning**

- In computation, the process is completed on its own
  - No interaction
    - The Turing machine automatically works according to programmed rules
  - A finite process
- In *learning*, there are two agents (teacher and learner)
  - Interaction between agents should be considered
    - A learning protocol between a teacher and a learner is essentially needed
  - An infinite process

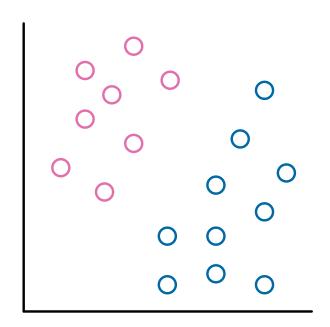
# Formalize Learning in Computational Manner

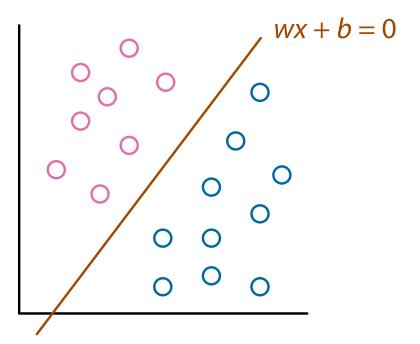
- 1. What are targets of learning?
- 2. How to represent targets and hypotheses?
- 3. How are data provided to a learner?
- 4. How does the learner work?
- 5. When can we say that the learner correctly learns the target?

## **Learning of Binary Classifier**



## **Learning of Binary Classifier**





#### **Example: Perceptron** (by F. Rosenblatt, 1958)

- **Learning target**: two subsets  $F, G \subseteq \mathbb{R}^d$  s.t.  $F \cap G = \emptyset$ 
  - Assumption: F and G are linearly separable
    - There exists a function (classifier)  $f_*(x) = w_*x + b$  s.t.  $f_*(x) > 0 \quad \forall x \in F, \qquad f_*(x) < 0 \quad \forall x \in G$

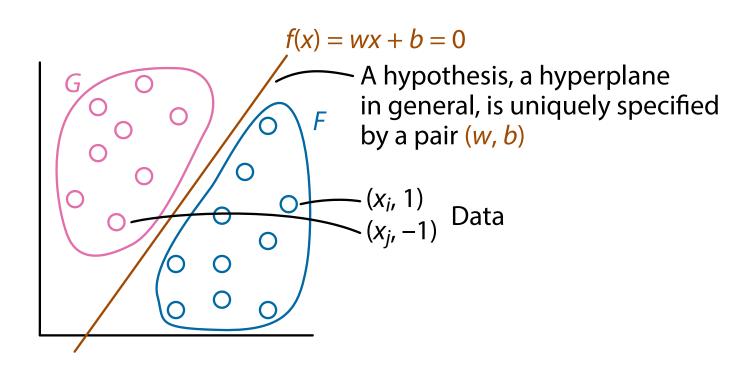
#### **Example: Perceptron** (by F. Rosenblatt, 1958)

- **Learning target**: two subsets  $F, G \subseteq \mathbb{R}^d$  s.t.  $F \cap G = \emptyset$ 
  - Assumption: F and G are linearly separable
    - There exists a function (classifier)  $f_*(x) = w_*x + b$  s.t.  $f_*(x) > 0 \quad \forall x \in F, \quad f_*(x) < 0 \quad \forall x \in G$
- **Hypotheses**: hyperplanes on  $\mathbb{R}^d$ 
  - If we consider a linear equation f(x) = wx + b, each line can be uniquely specified by a pair of two parameters (w, b) (hypothesis)

#### **Example: Perceptron** (by F. Rosenblatt, 1958)

- **Learning target**: two subsets  $F, G \subseteq \mathbb{R}^d$  s.t.  $F \cap G = \emptyset$ 
  - Assumption: F and G are linearly separable
    - There exists a function (classifier)  $f_*(x) = w_*x + b$  s.t.  $f_*(x) > 0 \quad \forall x \in F, \qquad f_*(x) < 0 \quad \forall x \in G$
- **Hypotheses**: hyperplanes on  $\mathbb{R}^d$ 
  - If we consider a linear equation f(x) = wx + b, each line can be uniquely specified by a pair of two parameters (w, b) (hypothesis)
- **Data**: a sequence of pairs  $(x_1, y_1), (x_2, y_2), ...$ 
  - $(x_i, y_i)$ : (a real-valued vector in  $\mathbb{R}^d$ , a label)
  - $x_i \in F \cup G$ ,  $y_i \in \{1, -1\}$ , and  $y_i = 1$   $(y_i = -1)$  if  $x_i \in F$   $(x_i \in G)$

## **Learning Model for Perceptron**



## **Learning Procedure of Perceptron**

1.  $w \leftarrow 0$ ,  $b \leftarrow 0$  (or a small random value)

// initialization

- 2. for i = 1, 2, 3, ... do
- 3. Receive *i*-th pair  $(x_i, y_i)$
- 4. Compute  $a = \sum_{j=1}^{d} w^j x_i^j + b^j$
- 5. if  $y_i \cdot a < 0$  then
- 6.  $w \leftarrow w + y_i x_i$
- 7.  $b \leftarrow b + y_i$
- 8. end if
- 9. end for

- //  $x_i$  is misclassified
- // update the weight
  - // update the bias

#### **Correctness of Perceptron**

- It is guaranteed that a perceptron always converges to a correct classifier
  - A correct classifier is a function f s.t.

$$f(x) > 0 \quad \forall x \in F,$$
  
 $f(x) < 0 \quad \forall x \in G$ 

- The convergence theorem
- Note: there are (infinitely) many functions that correctly classify F and G
  - A perceptron converges to one of them

## **Summary: Perceptron**

Target	Two disjoint subsets of $\mathbb{R}^d$
Representation	Two parameters $(w, b)$ of linear
	equation $f(x) = wx + b$
Data	Real vectors from target subsets
Algorithm	Perceptron
Correctness	Convergence theorem

#### **Example 2: Maximum Likelihood Estimation**

• Estimate the probability of a coin being a head in a toss

Target	Bernoulli distribution
Representation	Parameter (probability) $\it p$
Data	Sampling
Algorithm	Maximum Likelihood Estimation
	$\hat{p} = k/n$
Correctness	Consistency

## **Basic Definitions of Learning**

- Target: a classifier  $f_*: X \to \{0, 1\}$ 
  - A class  $\mathcal C$  of classifiers is usually pre-determined
  - Each target can be viewed as the set  $F_* = \{a \in X \mid f_*(a) = 1\}$ 
    - F<sub>\*</sub> is called a concept

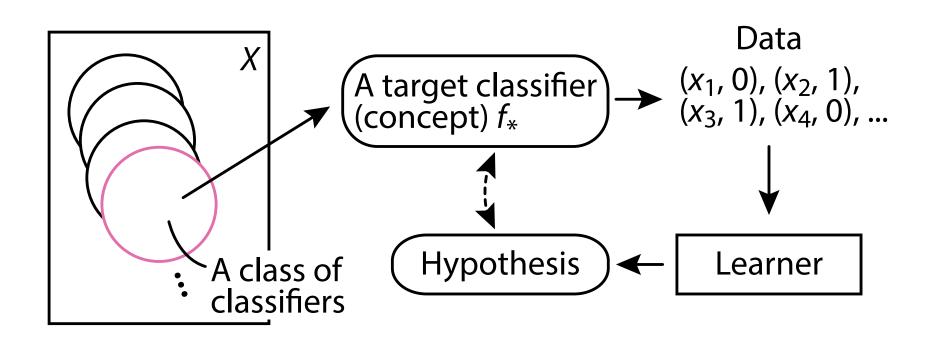
## **Basic Definitions of Learning**

- Target: a classifier  $f_*: X \to \{0, 1\}$ 
  - A class  $\mathcal C$  of classifiers is usually pre-determined
  - Each target can be viewed as the set  $F_* = \{a \in X \mid f_*(a) = 1\}$ 
    - F<sub>\*</sub> is called a concept
- Hypothesis space:  $\mathcal{R}$ 
  - Each hypothesis  $H \in \mathcal{R}$  represents a classifier
  - $\mathcal{R} \subseteq \Sigma^*$  usually holds ( $\Sigma^*$  is the set of finite strings)

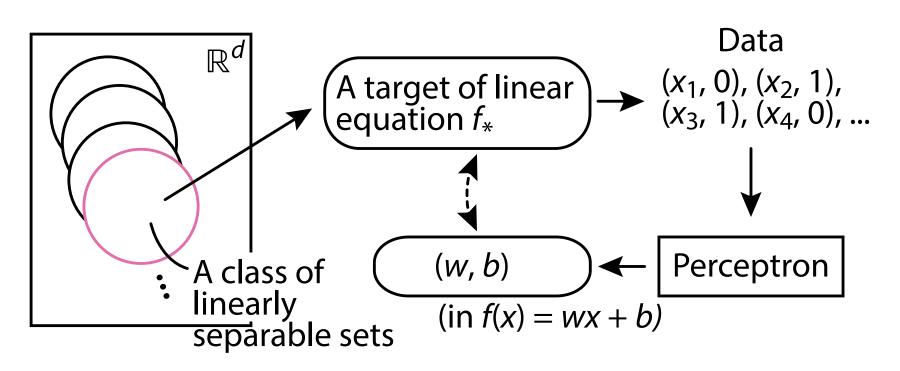
## **Basic Definitions of Learning**

- Target: a classifier  $f_*: X \to \{0, 1\}$ 
  - A class  $\mathcal{C}$  of classifiers is usually pre-determined
  - Each target can be viewed as the set  $F_* = \{a \in X \mid f_*(a) = 1\}$ 
    - F<sub>\*</sub> is called a concept
- Hypothesis space:  $\mathcal{R}$ 
  - Each hypothesis  $H \in \mathcal{R}$  represents a classifier
  - $\mathcal{R} \subseteq \Sigma^*$  usually holds ( $\Sigma^*$  is the set of finite strings)
- Data: Example  $(a, f_*(a))$ 
  - $-a \in X$
  - An example (a, 1) is called positive, (a, 0) is called negative

#### **Learning Model**



## Learning Model (e.g. Perceptron)



# Gold's Learning Model (Identification in the Limit)

- Gold gave the first basic learning model, called "Identification in the limit"
  - E. M. Gold, Language identification in the limit, Information and Control, 10(5), 447–474, 1967
- This model was originally introduced to analyze the learnability of formal languages
  - His motivation was to model infant's learning process of natural languages

#### Language Identification in the Limit

E MARK GOLD\*

The RAND Corporation

Language learnability has been investigated. This refers to the following situation: A class of possible languages is specified, together with a method of presenting information to the learner about an unknown language, which is to be chosen from the class. The question is now asked, "Is the information sufficient to determine which of the possible languages is the unknown language?" Many definitions of learnability are possible, but only the following is considered here: Time is quantized and has a finite starting time. At each time the

#### **Formal Languages**

- Alphabet  $\Sigma$ : a nonempty finite set
  - Each element  $a \in \Sigma$  is called a symbol
- Word  $w = a_1 a_2 \dots a_n$ : a finite sequence of symbols
  - Null word  $\varepsilon$ , whose length is 0
- The set of words  $\Sigma^*$  (with  $\varepsilon$ ) and  $\Sigma^+$  (without  $\varepsilon$ )

$$\Sigma^* = \{ a_1 a_2 \dots a_n \mid a_i \in \Sigma, n \ge 0 \}$$
  
$$\Sigma^+ = \{ a_1 a_2 \dots a_n \mid a_i \in \Sigma, n \ge 1 \} = \Sigma^* \setminus \{ \varepsilon \}$$

• Formal language: a subset of  $\Sigma^*$ 

#### Representation of Languages

- We connect syntax and semantics using a mapping f
- For a hypothesis  $H \in \mathcal{R}$ , f(H, w) is 0 or 1 for  $w \in \Sigma^*$ 
  - *H* is a program of a classifier
  - w is a (binary) code of the input to H
- $L(H) = \{ w \in \Sigma^* \mid f(H, w) = 1 \}$
- $\mathcal{R}$  is usually a recursively enumerable set
  - There is an algorithm that enumerates all elements of  ${\mathcal R}$
  - $\mathcal{R}$  is often identified with  $\mathbb{N}$ 
    - Each natural number encodes a classifier (hypothesis)

## **Setting of Gold's Learning Model**

- A class of languages  $\mathcal{C} \subseteq \{A \mid A \subseteq \Sigma^*\}$  is chosen
- For a language  $L \in \mathcal{C}$ , an infinite sequence  $\sigma = (x_1, y_1), (x_2, y_2), ...$  is a complete presentation of L if
  - (i)  $\{x_1, x_2, ...\} = \Sigma^*$
  - (ii)  $y_i = 1 \iff x_i \in L \text{ for all } i$ 
    - $\sigma[i] = (x_1, y_1), ..., (x_i, y_i)$  (a prefix of  $\sigma$ )
- A learner is a procedure M that receives  $\sigma$  and generates an infinite sequence of hypotheses  $\gamma = H_1, H_2, ...$ 
  - M outputs  $H_i$  if it gets  $\sigma[i]$

#### **Identification in the Limit**

- If  $\gamma$  converges to some hypothesis H and H represents L, we say that M identifies L in the limit
- If M identifies any  $L \in \mathcal{C}$  in the limit, we say that M identifies  $\mathcal{C}$  in the limit

#### **Basic Strategy: Generate and Test**

- Input: a complete presentation  $\sigma$  of a language L
- Output:  $\gamma = H_1, H_2, ...$
- 1.  $i \leftarrow 1, S \leftarrow \emptyset$
- 2. Repeat
- 3.  $S \leftarrow S \cup \{(x_i, y_i)\}$
- 4.  $k \leftarrow \min\{j \in \mathbb{N} \mid L(H(j)) \text{ consistent with } S\}$
- 5. // H(j) is a hypothesis encoded by a natural number j
- 6.  $H_i \leftarrow H(k)$  and output  $H_i$
- 7.  $i \leftarrow i + 1$
- 8. until forever

## **Power of Generate and Test Strategy**

- For any class  $\mathcal C$  of languages, Generate and Test strategy identifies  $\mathcal C$  in the limit
  - That is, Generate and Test strategy identifies every language  $L \in \mathcal{C}$  in the limit
- Unfortunately, this strategy is very inefficient
  - More intelligent strategy can be designed for each learning target
  - One of the most important tasks in studies of machine learning!

#### **Learning from Positive Data**

- In many cases, in particular in data mining, we obtain only positive data
  - Imagine supervised vs unsupervised learning
- A positive presentation of a language  $L \in \mathcal{C}$  is an infinite sequence  $x_1, x_2, ...$  s.t.  $L = \{x_1, x_2, ...\}$
- If  $\gamma$  (an infinite sequence of hypotheses of a learner M) converges to a hypothesis H s.t. L(H) = L, we say that M identifies L in the limit from positive data

## **Limitation of Learning from Positive Data**

- Consider the following class  $\mathcal C$ 
  - (i) All finite languages are included in  ${\mathcal C}$
  - (ii) At least one infinite language is included in  ${\mathcal C}$ 
    - *C* is called superfinite
- Gold proved that a superfinite class cannot be learned from positive data
  - e.g.  $\Sigma = \{a\}$ ,  $\mathcal{C}$  contains all finite languages and  $\{a^n \mid n \geq 1\}$
- Although this fact shows a limitation, there still exist rich classes of interesting languages
  - For example, pattern language

#### References

- If you are interested in computational learning theory, the following books might be interesting:
  - 榊原康文, 横森貴, 小林聡, 計算論的学習, 培風館, 2001
  - S. Jain, D. N. Osherson, J. S. Royer, A. Sharma, Systems That Learn, A Bradford Book, 1999
- These books are not necessarily for this lecture